

# THE SYMBOLIC DYNAMICS OF MULTIDIMENSIONAL TILING SYSTEMS

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ABSTRACT. We prove a multidimensional version of the theorem [CGST] that every shift of finite type has a power that can be realized as the same power of a tiling system. We also show that the set of entropies of tiling systems equals the set of entropies of shifts of finite type.

## 1. INTRODUCTION

In [CGST], E. Coven, W. Geller, S. Silberger, and W. Thurston consider a class of subshifts of  $\mathcal{A}^{\mathbb{Z}}$ , called tiles, defined by finite subsets of the integers. These subshifts are called tiling systems. The main result in [CGST] shows that given a one-dimensional shift of finite type, there is a tiling system and a positive integer  $m$ , such that the  $m$ -th powers of the shift of finite type and of the tiling system are topologically conjugate. It is also shown that every tiling system is sofic. It then follows that the set of entropies of tiling systems is equal to the set of entropies of shifts of finite type. In this paper we extend the main result to higher dimensions. In higher dimensions it is not known whether the set of entropies of shifts of finite type and of sofic shifts are equal, but we show that the set of entropies of tiling systems equals the set of entropies of shifts of finite type (Theorem 4.1).

We begin with some definitions. For a thorough introduction to symbolic dynamical systems, see either [K] or [LM]. Let  $\mathcal{A}$  be a finite alphabet, and consider the compact metric space  $\mathcal{A}^{\mathbb{Z}^d}$ . For  $x \in \mathcal{A}^{\mathbb{Z}^d}$  and  $w \in \mathbb{Z}^d$ , we will write  $x_w$  for  $x(w)$ . Let  $\sigma : \mathbb{Z}^d \rightarrow \text{Homeo}(\mathcal{A}^{\mathbb{Z}^d})$  be the continuous  $\mathbb{Z}^d$ -action (the homomorphism from  $\mathbb{Z}^d$  to the group of homeomorphisms of  $\mathcal{A}^{\mathbb{Z}^d}$ ) defined by

$$(\sigma_v(x))_w = x_{w+v}$$

for all  $v, w \in \mathbb{Z}^d$  and all  $x \in \mathcal{A}^{\mathbb{Z}^d}$ . Thus  $\sigma$  is defined by  $d$  commuting homeomorphisms  $\sigma_{e_i}$ ,  $i = 1, 2, \dots, d$ , where  $e_i$  is the  $d$ -tuple consisting of all 0's except for a 1 in the  $i$ -th position. We call  $\sigma$  the  $d$ -dimensional shift. For an integer  $m$  we define the  $m$ -th power,  $\sigma^m$ , of  $\sigma$  to be the  $\mathbb{Z}^d$ -action given by the composition homomorphism from  $\mathbb{Z}^d$  to  $\text{Homeo}(\mathcal{A}^{\mathbb{Z}^d})$ , which is defined by  $e_i \mapsto me_i \mapsto \sigma_{me_i}$  i.e.,

$$(\sigma_v^m(x))_w = x_{w+mv}.$$

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A compact (closed) set  $X \subseteq \mathcal{A}^{\mathbb{Z}^d}$  is said to be  $\sigma$ -invariant if for all  $v \in \mathbb{Z}^d$ ,  $\sigma_v(X) \subseteq X$ , or equivalently, if  $\sigma_{e_i}(X) \subseteq X$  for all  $i = 1, \dots, d$ . If  $X$  is a closed,  $\sigma$ -invariant subset of  $\mathcal{A}^{\mathbb{Z}^d}$ , we say that  $(X, \sigma)$  is a subshift of  $(\mathcal{A}^{\mathbb{Z}^d}, \sigma)$ .

We define a  $d$ -dimensional tiling system in the following way. Let  $\mathcal{P} = \{P_1, P_2, \dots, P_K\}$  be a finite collection of finite subsets of  $\mathbb{Z}^d$ , called *prototiles*. Normalize the prototiles so that the lexicographically smallest index is the origin. A *tile* is a translate of a prototile, and a *tiling of  $\mathbb{Z}^d$  by  $\mathcal{P}$*  is an expression of  $\mathbb{Z}^d$  as a disjoint union of tiles,  $\mathbb{Z}^d = \bigcup (t_j + P_{k_j})$ . Corresponding to a tiling is a point  $x \in \mathcal{A}^{\mathbb{Z}^d}$ ,  $\mathcal{A} = \{1, 2, \dots, K\}$ , where  $x_v = r$  if  $v$  lies in a tile that is a translate of  $P_r$ . Formally, for every  $x \in \mathcal{A}^{\mathbb{Z}^d}$  and for every  $v \in \mathbb{Z}^d$ ,  $x_v = r$  if and only if  $v \in P_{k_j} + t_j$  and  $k_j = r$ . Let  $T(\mathcal{P})$  denote the set of all points in  $\mathcal{A}^{\mathbb{Z}^d}$  corresponding to tilings of  $\mathbb{Z}^d$  by  $\mathcal{P}$ . The set  $T(\mathcal{P})$  is a  $\sigma$ -invariant, closed subset of the full shift  $\mathcal{A}^{\mathbb{Z}^d}$ , i.e., a subshift of  $\mathcal{A}^{\mathbb{Z}^d}$ . We call  $(T(\mathcal{P}), \sigma)$  a *tiling system*.

A  $d$ -dimensional shift of finite type is a subshift  $(X, \sigma)$  of  $(\mathcal{A}^{\mathbb{Z}^d}, \sigma)$  defined by a finite list of allowable  $d$ -dimensional cubes (see [LM], p. 467). By moving to a higher block presentation if necessary, a  $d$ -dimensional shift of finite type can be described as a closed, shift-invariant subspace  $\Sigma$  of  $\mathcal{A}^{\mathbb{Z}^d}$ , given by  $d$  directed, labeled graphs,  $G_1, \dots, G_d$ . For each  $1 \leq i \leq d$ , the mapping from the edges in  $G_i$  to the symbols in  $\mathcal{A}$  is one-to-one. A  $d$ -dimensional sequence  $x = (x_v)_{v \in \mathbb{Z}^d} \in \mathcal{A}^{\mathbb{Z}^d}$  is in  $\Sigma$  if and only if for each  $v \in \mathbb{Z}^d$  and for each  $i = 1, 2, \dots, d$ , the initial vertex of  $x_{v+e_i}$  in the graph  $G_i$  is the terminal vertex of  $x_v$  in  $G_i$ .

A  $d$ -dimensional sofic shift is a factor of a  $d$ -dimensional shift of finite type. The factor map is a block map that, without loss of generality, is defined on  $d$ -dimensional cubes. Again, by going to a higher block representation if necessary, every  $d$ -dimensional sofic shift is an image of a shift of finite type through a  $(1 \times 1 \times \dots \times 1)$ -block map. Unfortunately, no good graph representations of multidimensional sofic shifts are known. The straightforward extension of the one-dimensional case, changing the labels in the edge representation of the shift of finite type, does not necessarily represent the sofic shift (see Example 2.1).

For ease of notation, we state and prove the main theorem only in two dimensions. An outline of the proof for higher dimensions is included at the end of the paper. In two dimensions, the graphs representing a given shift of finite type will be denoted by  $G_H$  and  $G_V$  (for horizontal and vertical moves), instead of  $G_1$  and  $G_2$ . Similarly, we denote  $\sigma_{e_1}$  and  $\sigma_{e_2}$  by  $\sigma_H$  and  $\sigma_V$ .

**Theorem 1.1.** *Let  $(\Sigma, \sigma)$  be a two-dimensional shift of finite type. Then there is a positive integer  $m$  and a two-dimensional tiling system  $(T(\mathcal{P}), \sigma)$  such that*

1.  $T(\mathcal{P}) = \bigcup_{0 \leq i, j \leq m-1} T_{i,j}$ , where each  $T_{i,j}$  is a closed subset of  $T(\mathcal{P})$  and the collection  $\{T_{i,j}\}$  is permuted by  $\sigma$ :

$$\sigma_{(k,l)}(T_{i,j}) = T_{i+k,j+l},$$

where subscript addition is modulo  $m$ .

2.  $(\Sigma, \sigma^m)$  is topologically conjugate to  $(T_{i,j}, \sigma^m)$  for every  $(i, j)$ .

Using this theorem, we show that the set of topological entropies of multidimensional tiling systems equals the set of topological entropies of multidimensional shifts of finite type (Theorem 4.1). In one dimension, both sets are equal to the set of logarithms of roots of Perron numbers [LM] and to the set of entropies of sofic shifts. On the other hand, little is known about the set of topological entropies of multidimensional shifts of finite type and multidimensional sofic shifts.

The outline of the paper is as follows. In Section 2, we start with an example which illustrates some of the difficulties in working in more than one dimension and conclude with a proof that multidimensional tiling systems are sofic. In Section 3, we show that every shift of finite type can be realized as a power of a tiling system (Theorem 3.1). Although this result is not difficult, it helps to introduce the more difficult constructions needed for the proof of the main theorem. Section 4 contains the proof of Theorem 1.1, including a description of the prototiles of the two-dimensional tiling system. The modifications needed for the proof of Theorem 1.1 in higher dimensions are found in Section 5.

## 2. MULTIDIMENSIONAL TILING AND SOFIC SYSTEMS

In this section, we show that every tiling system is sofic. The one-dimensional case is proved in [CGST] using “subscripted tiling systems” and the “drop the subscripts” map. The proof for the multidimensional case can be done similarly. The idea of “dropping the subscripts” and the difficulty in presenting two-dimensional sofic shifts with graphs are illustrated in the example below.

**Example 2.1.** Let  $\mathcal{A} = \{a, b\}$ . Consider  $\Sigma \subseteq \mathcal{A}^{\mathbb{Z}^2}$ , a two-dimensional shift of finite type such that in every point  $x \in \Sigma$ , every appearance of  $b$  is surrounded by  $a$ 's. Then  $\Sigma$  is given by the following set of allowable  $2 \times 2$  blocks:

$$p := \begin{array}{cc} a & a \\ a & a \end{array} \quad q_1 := \begin{array}{cc} b & a \\ a & a \end{array} \quad q_2 := \begin{array}{cc} a & b \\ a & a \end{array} \quad q_3 := \begin{array}{cc} a & a \\ b & a \end{array} \quad q_4 := \begin{array}{cc} a & a \\ a & b \end{array}$$

By going to the  $(2 \times 2)$ -block presentation,  $\Sigma$  is topologically conjugate to  $\hat{\Sigma} \subseteq \{p, q_1, q_2, q_3, q_4\}^{\mathbb{Z}^2}$ , where  $\{p, q_1, q_2, q_3, q_4\}$  is the set of five allowable  $(2 \times 2)$ -blocks in  $\Sigma$ . The graphs representing  $\hat{\Sigma}$  are given in Figure 1.

Now consider the sofic system  $S$  obtained as a factor of  $\hat{\Sigma}$  using the  $(1 \times 1)$ -block map  $p \mapsto p$  and  $q_i \mapsto q$  for  $i = 1, \dots, 4$ . In [CGST], this map is called the “drop the subscripts” map.

What do the points in  $S$  look like? We can answer this by first looking at the points in  $\hat{\Sigma}$ . Note that if the symbol  $q_1$  appears in a point in  $\hat{\Sigma}$ , then it must be preceded horizontally by  $q_2$ . The symbol  $q_2$  must be followed vertically by  $q_4$ , and  $q_4$  must be followed horizontally by  $q_3$ . Hence, every appearance of the  $(1 \times 1)$ -block  $q_1$  occurs in an appearance of the  $(2 \times 2)$ -block  $\begin{array}{cc} q_4 & q_3 \\ q_2 & q_1 \end{array}$ . We come to the same conclusion if we consider the appearance of any other  $q_i$  ( $i = 2, 3, 4$ ). Thus every point in  $\hat{\Sigma}$  can be obtained by “concatenating”  $p$ 's and the block  $\begin{array}{cc} q_4 & q_3 \\ q_2 & q_1 \end{array}$  in a  $\mathbb{Z}^2$ -array. It then follows that every point in  $S$  can be obtained by concatenating  $p$ 's and the block  $\begin{array}{cc} q & q \\ q & q \end{array}$  in

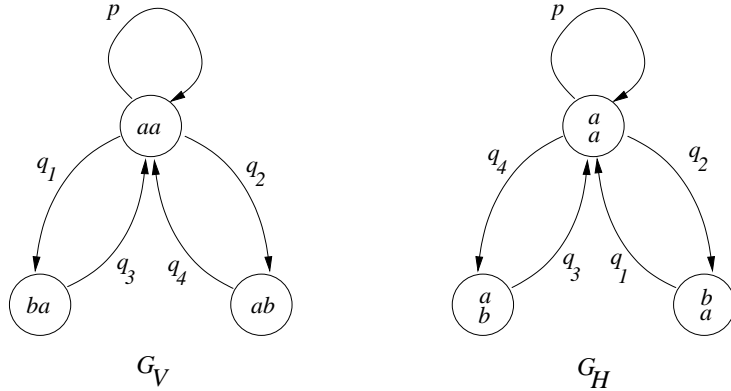
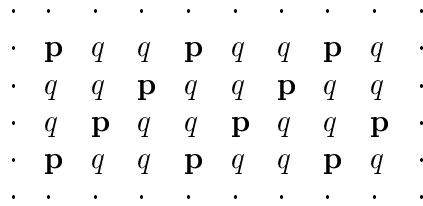


FIGURE 1. Graphs  $G_V$  and  $G_H$  representing the SFT  $\hat{\Sigma}$ .

a  $\mathbb{Z}^2$ -array. So  $S$  is the tiling system defined by two prototiles, a  $p$ -tile,  $p = \{(0, 0)\}$  and a  $q$ -tile,  $q = \{(0, 0), (0, 1), (1, 0), (1, 1)\}$ .

Note that  $S$  is *not* represented by the graphs obtained by dropping the subscripts of the edge labels of the graphs  $G_H$  and  $G_V$  in Figure 1. That presentation allows the following point, which is not in  $S$  ( $p$ 's are **bold** to make them easier to distinguish from  $q$ 's).



This example illustrates another difference between the “drop the subscripts” map in one-dimensional tiling systems and in multidimensional tiling systems. In one dimension, this map is finite-to-one and thus entropy-preserving. Our example shows that in more than one dimension, it can be infinite-to-one. The fixed point of all  $q$ 's is in  $S$  and has infinite number of preimages in  $\hat{\Sigma}$ . In fact,  $S$  cannot be an image of a shift of finite type by a finite-to-one map [Q]. However, Theorem 4.1 shows that the “drop the subscripts” map is entropy-preserving.

The following theorem shows that in all dimensions, tiling systems are sofic, but that the converse does not hold.

**Theorem 2.2.** *Every multidimensional tiling system is sofic. There is a two-dimensional shift of finite type that is not topologically conjugate to a tiling system.*

*Proof.* As mentioned before, the proof is essentially the same as the proof of the one-dimensional case [CGST]. We present an outline for two dimensions.

Let  $T(\mathcal{P})$  be a two-dimensional tiling system with  $\mathcal{P} = \{P_1, P_2, \dots, P_K\}$ . Consider the closed, shift-invariant subset  $\tilde{X}$  of  $\{k_{i,j} : 1 \leq k \leq K \text{ and } (i, j) \in P_k\}^{\mathbb{Z}^2}$ , defined by  $\tilde{x} \in \tilde{X}$  if and only if there is a tiling of  $\mathbb{Z}^2$  by  $\mathcal{P}$  such that for every  $(r, s) \in \mathbb{Z}^2$ ,  $\tilde{x}_{(r,s)} = k_{i,j}$ , where  $(r, s)$  is in the  $(i, j)$ -th location in  $P_k$ . (This is the two-dimensional version

of the one-dimensional “subscripted tiling system.”) As in the one-dimensional case,  $\tilde{X}$  is a shift of finite type, and the  $(1 \times 1)$ -block map  $k_{i,j} \mapsto k$  maps  $\tilde{X}$  onto  $T(\mathcal{P})$ . (This is the two-dimensional version of the “drop the subscripts” map.) Therefore  $T(\mathcal{P})$  is a sofic system.

Consider the shift of finite type  $\Sigma \subseteq \{0, 1\}^{\mathbb{Z}^2}$ , having exactly two points,  $x$  and  $y$ , such that  $\sigma_H(x) = y$ ,  $\sigma_H(y) = x$ ,  $\sigma_V(x) = x$ , and  $\sigma_V(y) = y$ . Then  $x_{2k,j} = 0 = y_{2k+1,j}$  and  $x_{2k+1,j} = 1 = y_{2k,j}$ , or the same is true with  $x$  and  $y$  interchanged. Notice that, except for interchanging the symbols, the only system topologically conjugate to  $\Sigma$  is  $\Sigma$  itself. Suppose that  $T$  is a tiling system that contains two points  $x_T$  and  $y_T$  such that  $\sigma_H(x_T) = y_T$ ,  $\sigma_H(y_T) = x_T$ ,  $\sigma_V(x_T) = x_T$ , and  $\sigma_V(y_T) = y_T$ . These two points are obtained from a set of two prototiles, a “0-tile” and a “1-tile.” The 0-tile can cover all points  $(i, j) \in \mathbb{Z}^2$  with  $i$  even and therefore it tiles  $\mathbb{Z}^2$ . Thus  $T$  has to contain the fixed point of all 0’s, and it cannot be conjugate to  $\Sigma$ .  $\square$

### 3. CONSTRUCTION OF THE TWO-DIMENSIONAL TILING SYSTEMS

From Theorem 2.2, there is a two-dimensional shift of finite type that is not a tiling system. In this section, we show that every two-dimensional shift of finite type is conjugate to a power of a subsystem of a tiling system. The proof of the following theorem also introduces the notion of marking the edges of a tile to reflect the transition behavior of the shift of finite type. This idea is modified in the proof of Theorem 1.1.

**Theorem 3.1.** *Let  $(\Sigma, \sigma)$  be a two-dimensional shift of finite type. Then there is a positive integer  $m$  and a two-dimensional tiling system  $(T(\mathcal{P}), \sigma)$  such that*

1.  $T(\mathcal{P}) = \bigcup_{0 \leq i, j \leq m-1} T_{i,j}$ , where each  $T_{i,j}$  is a closed subset of  $T(\mathcal{P})$  and the collection  $\{T_{i,j}\}$  is permuted by  $\sigma$ :

$$\sigma_{(k,l)}(T_{i,j}) = T_{i+k,j+l},$$

where subscript addition is modulo  $m$ .

2.  $(\Sigma, \sigma)$  is topologically conjugate to  $(T_{i,j}, \sigma^m)$  for every  $(i, j)$ .

*Proof.* The prototiles for the tiling system consist of squares, one associated to each symbol of  $\mathcal{A}$ , with “markings” on each side to govern how the squares can be put together.

For convenience in the construction, number the vertices in the horizontal graph  $G_H$ ,  $1, 2, 3, \dots, k_H$ , and in the vertical graph  $G_V$ ,  $1, 2, 3, \dots, k_V$ . Let  $m \geq 2 + \max(k_H, k_V)$ .

We assume that there are no parallel edges in  $G_V$  and in  $G_H$ . By moving to a higher block presentation if necessary, such graphs can be obtained for any shift of finite type.

The prototile associated to the symbol  $a$  is obtained as a modified  $m \times m$  square as follows: start with the  $m \times m$  square  $\{(i, j) : 0 \leq i, j \leq m-1\}$  in  $\mathbb{Z}^2$ . Above the north edge of the square, add the points  $(1, m), (2, m), \dots, (t_V(a), m)$ , where  $t_V(a)$  is the number associated to the terminal vertex of edge  $a$  in the vertical edge shift. To

the right of the east edge, add the points  $(m, m-2), (m, m-3), \dots, (m, t_H(a))$ , where  $t_H(a)$  is the number associated to the terminal vertex of edge  $a$  in the horizontal edge shift. From the south face, remove the points  $(1, 0), (2, 0), \dots, (i_V(a), 0)$ , where  $i_V(a)$  is the number associated to the initial vertex of edge  $a$  in the vertical edge shift. From the west face, remove the points  $(0, m-2), (0, m-3), \dots, (0, i_H(a))$ , where  $i_H(a)$  is the number associated to the initial vertex of edge  $a$  in the horizontal edge shift. We end up with an altered square that (when connected) looks like a jigsaw puzzle piece. This construction gives a one-to-one correspondence between the symbols in  $\mathcal{A}$  and the set of prototiles. For the tiling system considered in Example 2.1, the corresponding prototiles are depicted in Figure 2.

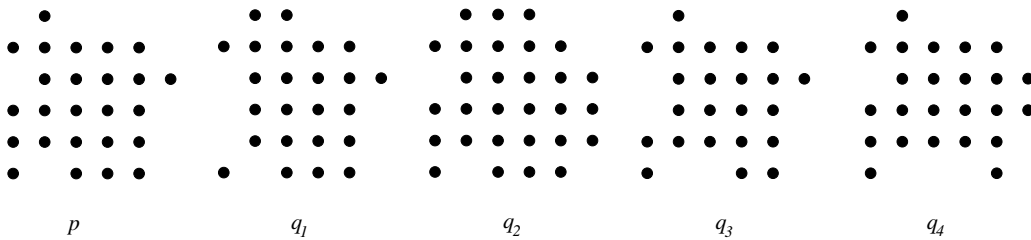


FIGURE 2. Prototiles associated to the SFT presented in Figure 1. The vertices  $\begin{smallmatrix} a \\ a \end{smallmatrix}, \begin{smallmatrix} a \\ b \end{smallmatrix}, \begin{smallmatrix} b \\ a \end{smallmatrix}$  in  $G_H$  and vertices  $aa, ba, ab$  in  $G_V$  are enumerated 1, 2, 3 respectively.

Let  $(T(\mathcal{P}), \sigma)$  be the resulting tiling system, and for  $0 \leq i, j \leq m-1$ , let  $T_{i,j}$  be the set of points in  $T(\mathcal{P})$  such that the southwest corners of the tiles appear at locations  $(mk + i, ml + j)$ . Let  $\varphi : T_{i,j} \rightarrow \Sigma$  be defined by  $(\varphi(t))_{k,l} = a$  if the tile whose southwest corner is at location  $(mk + i, ml + j)$  is the tile associated to symbol  $a$ . We leave it to the reader to verify that  $\varphi$  is a topological conjugacy of  $(T_{i,j}, \sigma^m)$  onto  $(\Sigma, \sigma)$ .  $\square$

#### 4. THE PROOF OF THEOREM 1.1

The prototiles necessary for the tiling system described in Theorem 1.1 must be more complicated than the ones described in Theorem 3.1. The authors in [CGST] use two kinds of prototiles, called *barbells* and *racks*. Here, we also use two kinds of prototiles, barbells and two-dimensional versions of racks.

##### Barbells and one-dimensional racks.

In [CGST], racks have three sections: a head, a center, and a tail. Barbells are used to fill the empty places in the center sections of the racks, and the number of ways this can be done equals the number of blocks of a fixed length in the given one-dimensional shift of finite type. The heads and tails of the racks are used to control which racks may sit side-by-side, i.e., the transition rules.

For the proof of Theorem 1.1 we also use barbells, and the center sections of the racks are used in the centers of the two-dimensional racks.

For fixed  $n$ , the center section of a rack has the form

$$\bullet^{3n+i} \xrightarrow{\leftarrow 2k \rightarrow} \bullet \xrightarrow{\leftarrow 2k \rightarrow} \bullet^{8n-4k-1-i}$$

where  $k = 1, \dots, n$  and  $i = 0, \dots, n$ . In this notation,  $\bullet$  is a point in the lattice of the two-dimensional integers,  $\bullet^s$  stands for  $\bullet \xrightarrow{\leftarrow s \rightarrow} \bullet$  and “ $\_$ ” denotes an empty space. We will denote this center section of a rack by  $r(k, i)$ .

Barbells are of the following form:

$$\bullet \bullet \xrightarrow{\leftarrow 1+2p \rightarrow} \bullet \bullet$$

where  $0 \leq p \leq 2n - 2$ . The gaps in the center,  $r(k, i)$ , of the rack may be filled in exactly  $k!$  ways by barbells. To see this, first imagine partitioning the set of empty places on the left side of the center into pairs of consecutive places. Then think of the placement of barbells in the rack as a one-to-one map from this partition onto the corresponding partition of the set of empty places on the right side of the center. There are exactly  $k!$  such maps. By varying  $i$  between 0 and  $n$ , we obtain  $n + 1$  different racks, and the centers of each of these racks can be filled in  $k!$  ways.

### Two-dimensional racks.

Beside barbells, for the two-dimensional case we use prototiles that are modified  $m \times m$  squares, which we call two-dimensional racks. The first modification uses the idea of side “markings” as described in the proof of Theorem 3.1. These markings control how the square tiles fit together on each side. An  $m \times m$  square tile intended to correspond to an  $(m \times m)$ -block in the original shift of finite type  $\Sigma$  has its sides marked in such a way that it can be adjacent to another  $m \times m$  square tile with marked edges if and only if the symbols on the sides of the corresponding  $(m \times m)$ -blocks in  $\Sigma$  can appear side-by-side. We call this first modification “side setting.”

The second modification involves leaving spaces in the center of the squares to be filled in a variety of ways by barbells. The spaces are chosen so that there is a one-to-one correspondence between the ways of filling these spaces with barbells and the  $(m \times m)$ -blocks in  $\Sigma$ . Such a correspondence leads naturally to a block map conjugacy between the tiling system we construct and the original shift of finite type  $\Sigma$ . This second modification is called “center setting.”

#### Step 1. Side setting.

Recall that the shift of finite type  $\Sigma$  is described by two edge graphs,  $G_H$  and  $G_V$ . Thus, symbol  $a$  can be followed by symbol  $b$  in the horizontal or vertical direction if the terminal vertex of  $a$  equals the initial vertex of  $b$  in the appropriate graph.

Now consider the  $(m \times m)$ -blocks over alphabet  $\mathcal{A}$  which are allowed in  $\Sigma$ . The block  $BB'$ , where both  $B$  and  $B'$  are allowable  $(m \times m)$ -blocks, is an allowable block in  $\Sigma$  if the right most column of  $B$  is “compatible” with the leftmost column of  $B'$ . That means: if the right most column of  $B$  is  $(a_1, \dots, a_m)^t$  and the leftmost column of  $B'$  is  $(a'_1, \dots, a'_m)^t$ , then  $a_i$  can be followed by  $a'_i$  in the horizontal direction for all  $i = 1, \dots, m$ . Let  $\mathcal{V}_H$  and  $\mathcal{V}_V$  be the sets of vertices of  $G_H$  and  $G_V$ , and  $\mathcal{V}_H^m$  and  $\mathcal{V}_V^m$  the set of  $m$ -tuples of vertices from these sets. Set  $\mathcal{V} = \mathcal{V}_H^m \cup \mathcal{V}_V^m$ .

Since the number of vertices is at most one more than the number of edges, we have  $|\mathcal{V}^m| \leq 2|\mathcal{A} + 1|^m < (2|\mathcal{A} + 1|)^m$ . Set  $s = 2|\mathcal{A} + 1|$ . There are at most  $|\mathcal{A}|^{m^2} < s^{m^2}$  allowable blocks of size  $m \times m$  in  $\Sigma$ . Because factorial growth is faster than exponential growth, we may choose  $n$  so that  $s^{13n} < (n!)$ . In fact, choose  $n$  so that  $|\mathcal{A}|^{(13n)^2} < s^{(13n)^2} < (n!)^n$ . Set  $m = 13n$ .

Let  $S_n$  be the group of permutations of  $n$  elements, and note that  $|S_n| = n!$ . Since  $|\mathcal{V}^m| < s^m < n!$ , there is an injection  $\pi$  from  $\mathcal{V}^m$  to  $S_n$ . Denote the permutation associated with the  $m$ -tuple  $\alpha \in \mathcal{V}^m$  by  $\pi_\alpha$ . We use  $\pi_\alpha$  to create a ‘‘comb-shaped marking’’ in the following way: let  $C(\pi_\alpha)$  be the subset of  $\mathbb{Z}^2$  consisting of the set of points  $\{(i, j) : 1 \leq i \leq n, 1 \leq j \leq \pi_\alpha(i)\}$ . In other words,  $C(\pi_\alpha)$  consists of  $n$  columns such that the  $i$ -th column has height  $\pi_\alpha(i)$ . We call the set  $\{(i, 1) : 1 \leq i \leq n\}$  the *base* of  $C(\pi_\alpha)$ .

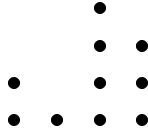


FIGURE 3.  $C(\pi_\alpha)$  for  $n = 4$  and  $\pi_\alpha = \begin{pmatrix} 1 & 2 & 3 & 4 \\ 2 & 1 & 4 & 3 \end{pmatrix}$ .

Figure 3 shows a comb facing north, which we will denote by  $C^\uparrow(\pi_\alpha)$ . We will also need combs facing south, east and west, which we will denote by  $C^\downarrow(\pi_\alpha)$ ,  $C^\rightarrow(\pi_\alpha)$ , and  $C^\leftarrow(\pi_\alpha)$ , respectively.



FIGURE 4.  $C^\leftarrow(\pi_\alpha)$  and  $\hat{C}^\rightarrow(\pi_\alpha)$  for  $\pi_\alpha = \begin{pmatrix} 1 & 2 & 3 & 4 \\ 2 & 1 & 4 & 3 \end{pmatrix}$ .

Associated to each comb  $C(\pi_\alpha)$  is its complement  $\hat{C}(\pi_\alpha) = \{(i, j) : 1 \leq i \leq n, 1 \leq j \leq n - \pi_\alpha(i)\}$  (see Figure 4). Note that  $C^\uparrow(\pi_\alpha) \cup \hat{C}^\downarrow(\pi_\beta)$  creates an  $n \times n$  square if and only if  $\pi_\alpha = \pi_\beta$ , that is if and only if the  $m$ -tuples of vertices  $\alpha$  and  $\beta$  are equal.

Given the boundary  $\mathcal{B}$  of an  $(m \times m)$ -block occurring in  $\Sigma$ , we use combs to ‘‘mark’’ the edges of an  $m \times m$  square tile as follows: Let  $\alpha_N, \alpha_S, \alpha_E, \alpha_W$  be the  $m$ -tuples of vertices associated to the one-dimensional  $m$ -blocks that make up the north, south, east, and west sides of  $\mathcal{B}$ . So  $\alpha_N, \alpha_S \in \mathcal{V}_V^m$  and  $\alpha_E, \alpha_W \in \mathcal{V}_H^m$ .

As described above,  $\pi_{\alpha_N}, \pi_{\alpha_S}, \pi_{\alpha_E}, \pi_{\alpha_W}$  are permutations and each has associated to it a comb shape. Suppose  $P \subseteq \mathbb{Z}^2$  denotes an  $m \times m$  square with corner vertices  $(0, 0)$ ,  $(m - 1, 0)$ ,  $(0, m - 1)$ , and  $(m - 1, m - 1)$ . The sides of  $P$  are altered by combs  $C^\uparrow(\pi_{\alpha_N})$ ,  $C^\rightarrow(\pi_{\alpha_E})$ ,  $\hat{C}^\leftarrow(\pi_{\alpha_W})$ , and  $\hat{C}^\downarrow(\pi_{\alpha_S})$ . We place  $C^\uparrow(\pi_{\alpha_N})$  so that its base is at points  $(6n, m), \dots, (7n - 1, m)$ . (Recall that  $m = 13n$ .) Place  $C^\rightarrow(\pi_{\alpha_E})$  so that its

base is at points  $(m, 6n), \dots, (m, 7n - 1)$ . Place  $\hat{C}^{\leftarrow}(\pi_{\alpha_W})$  and  $\hat{C}^{\downarrow}(\pi_{\alpha_S})$  so that the base of  $\hat{C}^{\leftarrow}(\pi_{\alpha_W})$  is at points  $(n - 1, 6n), \dots, (n - 1, 7n - 1)$ , and the base of  $\hat{C}^{\downarrow}(\pi_{\alpha_S})$  is at points  $(6n, n - 1), \dots, (7n - 1, n - 1)$ . This is illustrated in Figure 5 (but is not drawn to scale).

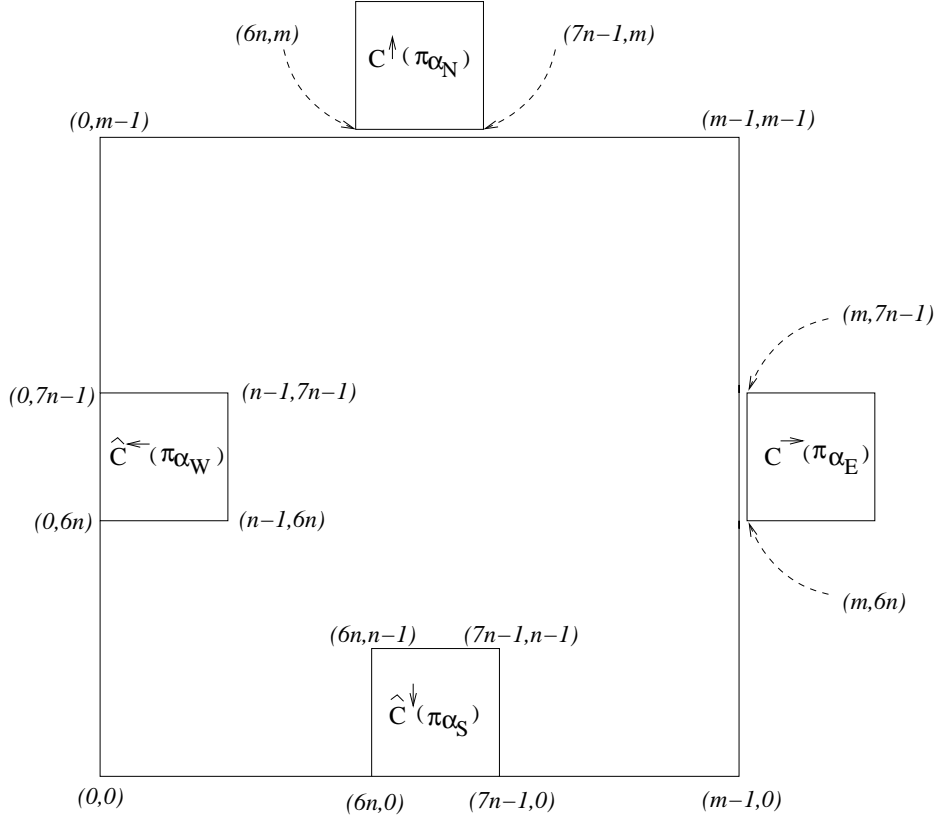


FIGURE 5. Side setting; sides of  $P$  are marked with combs,  $(m = 13n)$ .

We construct a square prototile, with edge combs as described above, for each of the boundaries  $\mathcal{B}$  that occur in  $(m \times m)$ -blocks in  $\Sigma$ . It is clear that these tiles can “fit together” horizontally (or vertically) if and only if the corresponding boundaries of  $(m \times m)$ -blocks in  $\Sigma$  can occur horizontally (or vertically) adjacent to one another for some point in  $\Sigma$ .

### Step 2. Center Setting

There may be many  $(m \times m)$ -blocks in  $\Sigma$  with the same boundary, and we make our second modification to reflect this. As in the one-dimensional construction, we introduce spaces into the central portion of the modified tile. This is done so that there are exactly as many ways of filling in the central spaces in the prototiles associated with boundary  $\mathcal{B}$  as there are  $(m \times m)$ -blocks in  $\Sigma$  with boundary  $\mathcal{B}$ . The tile  $P$ , once it is modified with “side setting” and “center setting”, is a two-dimensional rack.

Given a boundary  $\mathcal{B}$  of an  $(m \times m)$ -block occurring in  $\Sigma$ , let  $N_{\mathcal{B}}$  be the number of allowable  $(m \times m)$ -blocks in  $\Sigma$  with boundary  $\mathcal{B}$ . By our choice of  $n$ , it is clear that

$$N_{\mathcal{B}} \leq |\mathcal{A}|^{(13n)^2} < s^{(13n)^2} < (n!)^n,$$

and thus  $N_{\mathcal{B}}$  can be written uniquely (in base  $n!$ ) as follows:

$$N_{\mathcal{B}} = d_0 + d_1(n!) + d_2(n!)^2 + \cdots + d_{n-1}(n!)^{n-1},$$

where  $0 \leq d_j < n!$  for  $0 \leq j \leq n-1$ . As in [CGST], each  $d_j$  can be written uniquely as

$$d_j = c_{j,1}(1!) + c_{j,2}(2!) + \cdots + c_{j,n-1}(n-1!),$$

where  $0 \leq c_{j,k} \leq k$  and  $1 \leq k \leq n-1$ . Thus

$$N_{\mathcal{B}} = \sum_{j=0}^{n-1} d_j(n!)^j = \sum_{j=0}^{n-1} \sum_{k=1}^{n-1} c_{j,k}(k!)(n!)^j.$$

Begin with the prototile  $P$ , altered with combs corresponding to  $\mathcal{B}$  as illustrated in Figure 5. For each  $0 \leq j \leq n-1$  and each  $1 \leq k \leq n-1$ , form a two-dimensional rack  $P_i(j, k)$  obtained from  $P$  in the following way. Consider a rectangular section of size  $11n \times (j+1)$  such that the bottom  $j$  rows are filled with one-dimensional racks  $r(n, 0)$  and the top row with the rack  $r(k, i)$ . The two-dimensional rack  $P_i(j, k)$  is obtained from  $P$  by replacing the section with corners  $(n, n+1)$ ,  $(12n-1, n+1)$ ,  $(n, n+j+1)$ ,  $(12n-1, n+j+1)$  with this  $11n \times (j+1)$  rectangular shape.

Observe that by construction of the one-dimensional racks, the bottom  $j$  rows can be filled by barbells in exactly  $n!$  ways and the top row can be filled in exactly  $k!$  ways.

Thus, the modified prototile  $P_i(j, k)$  can have the empty spaces filled with barbells in exactly  $(k!)(n!)^j$  ways. By varying  $i$  between 1 and  $c_{j,k}$ , we obtain  $c_{j,k}$  prototiles whose empty spaces can be filled with barbells in  $(k!)(n!)^j$  ways.

So, the empty places in the racks  $P_1(j, k), \dots, P_{c_{j,k}}(j, k)$  can together be filled with barbells in  $c_{j,k}(k!)(n!)^j$  ways. Let  $\mathcal{P}(j, k) = \{P_i(j, k) : 1 \leq i \leq c_{j,k}\}$ . Then the set of two-dimensional racks that correspond to the boundary  $\mathcal{B}$  is  $\mathcal{P}_{\mathcal{B}} = \bigcup_{j=0}^{n-1} \bigcup_{k=1}^{n-1} \mathcal{P}(j, k)$ . Hence the empty spaces of the racks in  $\mathcal{P}_{\mathcal{B}}$  can be filled by barbells in exactly  $N_{\mathcal{B}}$  ways.

### Proof of Theorem 1.1.

Let  $\Sigma^m$  be the  $(m \times m)$ -power of  $\Sigma$ . As in the one-dimensional case,  $\Sigma^m$  is defined by  $(m \times m)$ -blocks that are allowable in  $\Sigma$ , i.e.,  $\Sigma^m \subseteq (A^{m^2})^{\mathbb{Z}^2}$ . Hence  $(\Sigma^m, \sigma)$  is conjugate to  $(\Sigma, \sigma^m)$  [LM]. We show that the tiling system constructed above with barbells and two-dimensional racks as prototiles is conjugate to  $(\Sigma^m, \sigma)$ .

First, observe that the barbells can be used to fill only the spaces in the centers of rack tiles. The barbells, just by themselves, cannot be used to tile any rectangular region because they have an odd number of empty spaces in the middle and they cover an even number of spaces. Clearly, they cannot be used to fill the empty spaces in the combs marking the edges of the rack tiles. Thus, the only way to tile the

plane with these tiles is for racks to sit next to each other with their comb markings compatible, and for barbells to fill the spaces inside the racks.

Let  $T$  be the tiling system with the collection of prototiles  $\mathcal{P}$  being the barbells and racks. Let  $T_{0,0}$  denote the collection of all tilings in  $T$  that consist of racks occurring with their southwest corners at lattice points  $\{(km, jm) : k, j \in \mathbb{Z}\}$ . Clearly  $T = \bigcup_{0 \leq i, j \leq m-1} T_{i,j}$ , where  $T_{i,j} = \sigma_{(i,j)} T_{0,0}$ . This gives statement (1) of Theorem 1.1.

To prove statement (2) of Theorem 1.1, let  $G_H^m$  be the directed graph with vertices  $\mathcal{V}_H^m$ , given by  $m$ -tuples of vertices in  $\mathcal{V}_H$ .

Let  $x \in \Sigma$ . For each  $x_{i,j}$ , let  $I(x_{i,j})$  and  $J(x_{i,j})$  denote the initial and terminal vertices of the edge in  $G_H$  with label  $x_{i,j}$ . For every rack  $P$  with the comb marking on its west side constructed using  $(I(x_{0,0}), \dots, I(x_{0,m-1})) \in \mathcal{V}_H$  and the comb marking on its east side constructed using  $(J(x_{m-1,0}), \dots, J(x_{m-1,m-1}))$ , put an edge from vertex  $(I(x_{0,0}), \dots, I(x_{0,m-1}))$  to vertex  $(J(x_{m-1,0}), \dots, J(x_{m-1,m-1}))$ . Label this edge  $P$ .

Similarly, let  $G_V^m$  be the directed graph with vertices  $\mathcal{V}_V^m$ , the  $m$ -tuples of vertices in  $\mathcal{V}_V$ . In this graph, the edges are determined by the north and south sides of the racks  $P$ .

Then  $G_H^m$  and  $G_V^m$  correspond exactly to the directed graphs of  $(\Sigma^m, \sigma)$ . Thus  $(T_{i,j}, \sigma^m)$  is topologically conjugate to the  $(m \times m)$ -power  $(\Sigma^m, \sigma)$ , which is topologically conjugate to  $(\Sigma, \sigma^m)$ , as desired.  $\square$

**Theorem 4.1.** [B]

$$\begin{aligned} \{h(\Sigma, \sigma) : \Sigma \text{ is a shift of finite type}\} &= \{h(T, \sigma) : T \text{ is a tiling system}\} \\ &\subseteq \{h(S, \sigma) : S \text{ is a sofic system}\}, \end{aligned}$$

where  $h(\cdot)$  denotes two-dimensional topological entropy.

*Proof.* From Theorem 1.1 and Proposition 13.1 in [S], we have that  $\{h(\Sigma, \sigma) : \Sigma \text{ is a shift of finite type}\} \subseteq \{h(T, \sigma) : T \text{ is a tiling system}\}$ . The inclusion  $\{h(T, \sigma) : T \text{ is a tiling system}\} \subseteq \{h(S, \sigma) : S \text{ is a sofic system}\}$  follows from Theorem 2.2. The statement of the Theorem will follow if we show that “drop the subscripts” map is entropy-preserving. Again we prove the statement in two dimensions only. The higher-dimensional case follows similarly.

Let  $T = T(\mathcal{P})$  be a two-dimensional tiling system defined by a set of prototiles  $\mathcal{P} = \{P_1, \dots, P_K\}$ . Let  $\hat{T}$  be the closed, shift-invariant subset of  $\{k_{i,j} : 1 \leq k \leq K \text{ and } (i, j) \in P_k\}^{\mathbb{Z}^2}$  described in the proof of Theorem 2.2, and let  $\delta : \hat{T} \rightarrow T$  be the  $(1 \times 1)$ -block “drop the subscripts” map. The prototiles that define  $\hat{T}$  are subscripted prototiles  $\hat{\mathcal{P}} = \{\hat{P}_1, \dots, \hat{P}_K\}$ . Set  $s = \max\{|i_1 - i_2|, |j_1 - j_2| : (i_1, j_1), (i_2, j_2) \in P_r, r = 1, \dots, K\}$ . Then  $\hat{T}$  is a shift of finite type defined by  $(s \times s)$ -blocks.

Consider an allowable  $((n + 2s) \times (n + 2s))$ -block  $B$  in  $T$  and  $((n + 2s) \times (n + 2s))$ -blocks  $\hat{B}_1, \hat{B}_2$  in  $\hat{T}$  such that  $\delta(\hat{B}_1) = \delta(\hat{B}_2) = B$ . Let  $C, \hat{C}_1$  and  $\hat{C}_2$  be the central  $n \times n$  portions of  $B, \hat{B}_1, \hat{B}_2$  respectively. We show that if  $\hat{B}_1$  and  $\hat{B}_2$  have equal boundaries of thickness  $s$ , then their central portions must be equal (in other words, the “drop the subscripts” map “has no diamonds” (see [LM],[K])).

Since  $\delta(\hat{B}_1) = \delta(\hat{B}_2) = B$ , it is clear that the entries in the central portions  $\hat{C}_1$  and  $\hat{C}_2$  can differ only in their subscripts. In other words, each entry must be in a

translate of the same prototile. If that translate intersects the  $s$ -thick boundary of either  $\hat{B}_i$ , then the subscripts clearly must be equal. If neither translated prototile intersects the  $s$ -thick boundary and yet the subscripts of the two entries are different, there must be two different ways to tile a portion of the centers of the  $\hat{B}_i$ 's with the translated prototile.

Let  $P \in \mathcal{P}$  and  $t_i, r_i \in \mathbb{Z}^2$  for  $i = 1, \dots, k$ . Let  $T_k = \bigcup_{i=1}^k (t_i + P)$  and  $R_k = \bigcup_{i=1}^k (r_i + P)$ , where the unions are disjoint.

Then

$$T_k = R_k \text{ implies } \{t_1, \dots, t_k\} = \{r_1, \dots, r_k\}. \quad (*)$$

We show  $(*)$  by induction on  $k$ . The case  $k = 1$  is trivial. Assume that  $T_k = R_k$ . As mentioned in the Introduction, the prototile  $P$  is chosen so that the lexicographically smallest point in  $P$  is  $(0, 0)$ , and so the smallest point in  $t_i + P$  is  $t_i$ . Thus, if  $v$  is the lexicographically smallest point in  $T_k = R_k$ ,  $v$  is in both  $\{t_1, \dots, t_k\}$  and  $\{r_1, \dots, r_k\}$ . Without loss of generality,  $v = t_k = r_k$ . But then  $T_{k-1} = R_{k-1}$  and  $(*)$  follows from the induction hypothesis.

Now let  $r_i \in \mathbb{Z}^2$  ( $i = 1, \dots, k$ ) be such that for some prototile  $P$ ,  $r_i + \hat{P} \subseteq \hat{B}_1$ . Similarly, let  $t_i \in \mathbb{Z}^2$  ( $i = 1, \dots, k'$ ) be such that  $t_i + \hat{P} \subseteq \hat{B}_2$ . Then  $r_i$  and  $t_i$  are either in the central portions  $\hat{C}_1$  and  $\hat{C}_2$  or in the boundaries with thickness  $s$  of  $\hat{B}_1$  and  $\hat{B}_2$ . But  $\hat{B}_1$  and  $\hat{B}_2$  have the same boundary and  $\delta(\hat{B}_1) = \delta(\hat{B}_2) = B$ . So we must have  $k = k'$  and  $\bigcup_{i=1}^k (r_i + \hat{P}) = \bigcup_{i=1}^k (t_i + \hat{P})$ , which by  $(*)$  implies  $\{r_1, \dots, r_k\} = \{t_1, \dots, t_k\}$ . Thus the central portions  $\hat{C}_1$  and  $\hat{C}_2$  are equal.

Now, let  $\hat{N}_n$  be the number of  $(n \times n)$ -blocks in  $\hat{T}$  ( $n > 2s$ ),  $\hat{K}$  be the number of symbols in  $\hat{T}$ , and  $\mathcal{N}_n$  the number of  $(n \times n)$ -blocks in  $T$ . The number of  $s$ -thick boundaries for an  $(n \times n)$ -block in  $\hat{T}$  is at most  $\hat{K}^{4s(n-s)}$ . Then the corollary follows from the inequality

$$\mathcal{N}_n \leq \hat{N}_n \leq \hat{K}^{4s(n-s)} \mathcal{N}_{n-2s}$$

and the fact that since  $s$  is fixed, the exponent of  $\hat{K}$  is linear in  $n$ .  $\square$

## 5. HIGHER DIMENSIONS

The construction of the two-dimensional racks described in the previous section can be extended inductively to  $d$ -dimensional racks. Given a boundary  $\mathcal{B}$  of an  $(m \times m \times \dots \times m)$ -block occurring in a  $d$ -dimensional shift of finite type  $\Sigma$ , let  $N_{\mathcal{B}}$  be the number of  $(m \times m \times \dots \times m)$ -blocks with boundary  $\mathcal{B}$ . The boundary  $\mathcal{B}$  consists of the  $(d-1)$ -dimensional faces of the block; there are at most  $|\mathcal{A}+1|^{m^{d-1}}$  possibilities for each of these faces. Choose  $n > 0$  large enough so that

$$N_{\mathcal{B}} \leq |\mathcal{A}|^{(13n)^d} < |d(\mathcal{A}+1)|^{(13n)^d} < (n!)^{(n^{d-1})}.$$

Thus we may write  $N_{\mathcal{B}}$  uniquely (in base  $(n!)^{n^{d-2}}$ ) as follows:

$$N_{\mathcal{B}} = b_0 + b_1((n!)^{n^{d-2}}) + b_2((n!)^{n^{d-2}})^2 + \dots + b_{n-1}((n!)^{n^{d-2}})^{n-1},$$

where  $0 \leq b_j < (n!)^{n^{d-2}}$  for  $0 \leq j \leq n-1$ . For each  $b_j$ , write

$$b_j = c_{j,1} + c_{j,2}((n!)^{n^{d-3}}) + \dots + c_{j,n}((n!)^{n^{d-3}})^{n-1},$$

where  $0 \leq c_{j,k} < (n!)^{n^{d-3}}$ . Thus

$$N_{\mathcal{B}} = \sum_{j=0}^{n-1} \sum_{k=1}^{n-1} c_{j,k} ((n!)^{d-3})^{k-1} ((n!)^{n^{d-2}})^j.$$

Set  $m = 13n$ .

In the construction of the prototiles at the previous stage, dimension  $d - 1$ , the spaces in the center sections of the  $(d - 1)$ -dimensional racks can be filled in with barbells in  $b$  ways, where  $b$  is any number between 0 and  $(n!)^{n^{d-2}}$ . These center sections have size  $11n \times (n \times \overset{\leftarrow{d-2}\rightarrow}{n \times \cdots \times n})$ . (Note the  $n$  at the previous stage may have been smaller; there is no difficulty in increasing it.) Thus, in  $d$  dimensions, for each term in the sum above, we construct a collection of  $(m \times m \times \cdots \times m)$ -block prototiles. For each  $j$ ,  $0 \leq j \leq n - 1$ , we replace a block of size  $11n \times (n \times \overset{\leftarrow{d-2}\rightarrow}{n \times \cdots \times n}) \times (j + 1)$  in the center portion of this prototile. The first  $j$  “sheets” of size  $11n \times (n \times \overset{\leftarrow{d-2}\rightarrow}{n \times \cdots \times n})$  are constructed as in the  $(d - 1)$ -dimensional case, each to be filled with barbells in exactly  $(n!)^{n^{d-2}}$  ways. So, all together, they can be filled in  $((n!)^{n^{d-2}})^j$  ways. The last “sheet” of size  $11n \times (n \times \overset{\leftarrow{d-2}\rightarrow}{n \times \cdots \times n})$  is constructed as in the  $(d - 1)$ -dimensional case, to be filled with barbells in exactly  $((n!)^{n^{d-3}})^{k-1}$  ways. By varying  $k$ , we obtain  $c_{j,k}$  different  $d$ -dimensional racks with center sections that can be filled by barbells in  $c_{j,k} ((n!)^{n^{d-3}})^{k-1} ((n!)^{n^{d-2}})^j$  ways.

We do this for each  $j$ ,  $0 \leq j \leq n - 1$ , and thus the centers of this collection of  $(m \times m \times \cdots \times m)$ -prototiles can be filled with barbells in exactly  $N_{\mathcal{B}}$  ways.

It remains only to mark the bounding faces of these prototiles with combs appropriate for  $\mathcal{B}$ . We note that, due to our choice of  $n$ , the number of possibilities for a  $(d - 1)$ -dimensional face of an  $(m \times m \times \cdots \times m)$ -block in  $\Sigma$  is bounded by  $d|\mathcal{A} + 1|^{(13n)^{d-1}} < (d(|\mathcal{A}| + 1))^{(13n)^{d-1}} < (n!)^{(n^{d-2})}$ . Thus there is an injection from the set of faces  $\bigcup_{i=1}^d \mathcal{V}_i^{(13n)^{d-1}}$  (where  $\mathcal{V}_i$  is the set of vertices in the  $i$ -th graph representing  $\Sigma$ ) into the set  $S_n^{d-2}$  (i.e., the set of  $(d - 2)$ -tuples of elements of  $S_n$ ). We will mark each  $(d - 1)$ -dimensional face of the prototile with the  $n^{d-2}$  combs associated with that face, each of which represents a permutation in  $S_n$ .

It is clear that, as in two dimensions, the barbells can be used to fill only the spaces in the centers of the  $d$ -dimensional racks. Thus the only way to tile the  $d$ -dimensional integer lattice with these tiles is for the  $d$ -dimensional racks to sit next to each other with their comb markings compatible, and for the barbells to fill the spaces inside the  $d$ -dimensional racks. The proof of the theorem in  $d$ -dimensions then follows analogously.

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